

# MIPS assembly

# Review

- We learned
  - `addi`,
  - `and`, `andi`, `or`, `ori`, `xor`, `xori`,
  - `nor`,
- An array is stored sequentially in the memory
- The instructions are also stored sequentially in the memory. Executing the code is to load then execute the instructions one by one, unless we encounter a branch condition.

# Shifts

- Shift instructions move all the bits in a word to the left or to the right
  - Shift left logical (sll) move all the bits to the left by the specified number of bits
    - `sll $t2, $t0, 2`
  - Shift right logical (srl) move all the bits to the right
    - `srl $t2, $t0, 2`
  - Filling the emptied bits with 0's
    - This includes srl with negative numbers (since you insert 0's to the left of the number, your number will be positive after the shift)

# Example 1

- Suppose register \$s0 (\$16) is 9<sub>ten</sub>

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	1

- What do we have in \$t2 (\$10) after `sll $t2, $s0, 4`

# Example 1

- Suppose register \$s0 (\$16) is  $9_{\text{ten}}$

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	1

– We have in \$t2 (\$10) after `sll $t2, $s0, 4`

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	1	0	0	0	0	

- The value is  $144_{\text{ten}} = 9_{\text{ten}} \times 2^4$
- In general, shifting left by  $i$  bits gives the same result as multiplying by  $2^i$

# Example 2

- Suppose register \$s0 (\$16) is 9<sub>ten</sub>

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	1

- What do we have in \$t2 (\$10) after `sll $t2, $s0, 28`

# Example 2

- Suppose register \$s0 (\$16) is  $9_{\text{ten}}$

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	1

– We have in \$t2 (\$10) after `sll $t2, $s0, 28`

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

- The value is NOT  $9_{\text{ten}} \times 2^{28}$  noting that the number is a signed number.
- Overflow happens this time

# Example 3

- Suppose register \$s0 (\$16) is 99<sub>ten</sub>

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	0	1	1

- What do we have in \$t2 (\$10) after `srl $t2, $s0, 4`



# Example 3

- Suppose register \$s0 (\$16) is  $99_{\text{ten}}$

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	0	1	1

- We have in \$t2 (\$10) after `srl $t2, $s0, 4`

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0

- The value is  $6_{\text{ten}} = 99_{\text{ten}} / 2^4$
- In general, shifting left by  $i$  bits gives the same result as dividing by  $2^i$

# Example 4

- Suppose register \$s0 (\$16) is  $-9_{\text{ten}}$

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1	1	1

- What do we have in \$t2 (\$10) after `srl $t2, $s0, 4`

# Example 4

- Suppose register \$s0 (\$16) is  $-9_{\text{ten}}$

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1	1	1

– We have in \$t2 (\$10) after srl \$t2, \$s0, 4

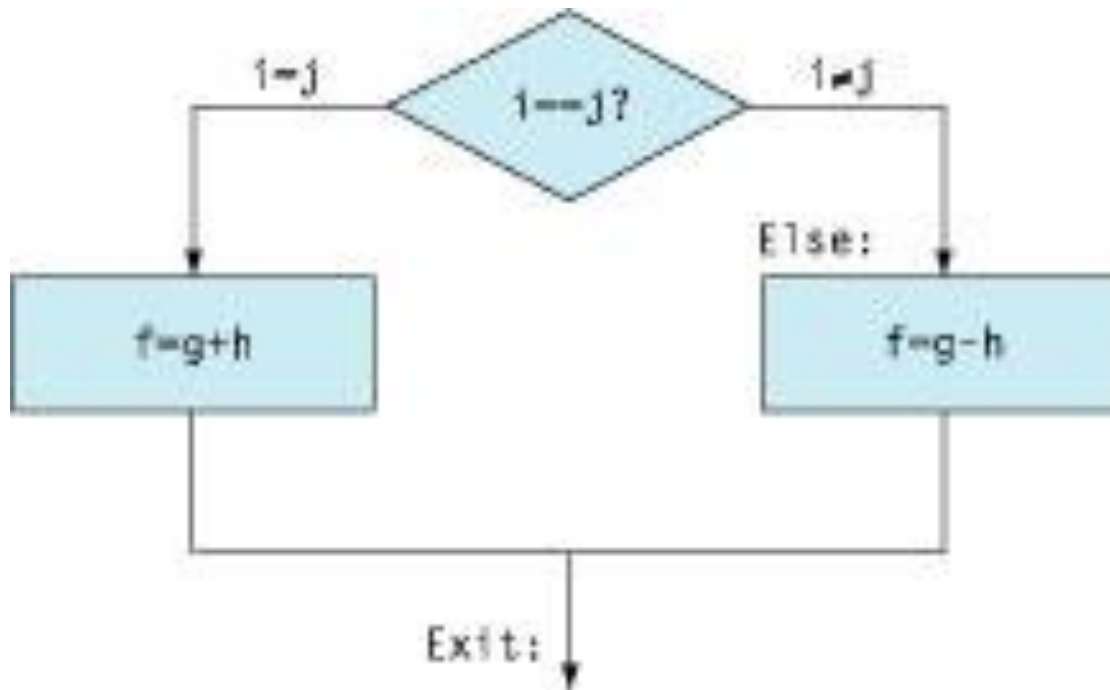
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	

- The value is NOT  $-9_{\text{ten}} / 2^4$  noting that the number is a signed number.
- Even though it's a negative number, 0's are filled in during shift

# Instructions for Making Decisions

- A distinctive feature of programs is that they can make different decisions based on the input data

```
if (i==j) f = g + h; else f = g - h;
```



# Instruction beq (branch if equal)

- To support decision making, MIPS has two conditional branch instructions, similar to an “if” statement with a goto

```
beq register1, register2, L1
```

- In C, it is equivalent to

```
if (register1 == register2)
    goto L1
```

- Note that L1 is a label and we are comparing values in register1 and register2

- **Label is an address of an instruction.**
  - Every address can be associated with a label, which is used by the assembly program to specify the address
  - Go to a label means that fetch that instruction from the memory and execute it.

# Instruction bne

- Similarly, bne (branch not equal) means go to the statement labeled with L1 if the value in register1 does not equal to the value in register2

```
bne register1, register2, L1
```

– Equivalent to

```
if (register1 != register2)  
    goto L1
```

# Instruction j (jump)

- MIPS has also an unconditional branch, equivalent to goto in C

j L1

– Jump to the instruction labeled with L1

# Compiling if-then-else

- Suppose variables  $f$ ,  $g$ ,  $h$ ,  $i$ , and  $j$  are in registers  $\$s0$  through  $\$s4$ , how to implement the following in MIPS?

```
if (i==j) f = g + h; else f = g - h;
```



# Compiling if-then-else

- Suppose variables f, g, h, i, and j are in registers \$s0 through \$s4, how to implement the following in MIPS?

```
if (i==j) f = g + h; else f = g - h;
```

```
    if (i != j)
        goto Else;
    f = g + h;
    goto Exit;
Else:
    f = g - h;
Exit:
```

# Compiling if-then-else

- Suppose variables f, g, h, i, and j are in registers \$s0 through \$s4, how to implement the following in MIPS?

```
if (i==j) f = g + h; else f = g - h;
```

```
    if ($s3 != $s4)
        goto Else;
    $s0 = $s1 + $s2;
    goto Exit;
```

```
Else:
```

```
    $s0 = $s1 - $s2;
```

```
Exit:
```

# MIPS Assembly for if-then-else

- Now it is straightforward to translate the C program into MIPS assembly

```
if (i==j) f = g + h; else f = g - h;
```

```
        bne $s3,$s4,Else;      #go to Else if i <> j
        add $s0, $s1, $s2     #f = g + h
        j    Exit;           #go to the end of the if-then-else block
Else:
        sub $s0, $s1, $s2     #f = g -h
Exit:
```

# Exercise 1

- Suppose \$t0 is storing 30, \$t1 is storing 20. After the following instructions, what will be the value in \$t2?

```
sub $t2, $t0, $t1
```

```
srl $t2, $t2, 2
```

```
ori $t2, $t2, 10
```

(a) 8

(b) 10

(c) 18

(d) None of the above.

# Exercise 2

- Suppose word array A stores 0,1,2,3,4,5,6,7,8,9, in this order. Assume the starting address of A is in \$s0. After the following instructions, what will be the value in \$t0?

```
addi $s0, $s0, 32
```

```
lw $t0, 4($s0)
```

```
andi $t0, $t0, 1
```

- (a) 0
- (b) 8
- (c) 9
- (d) None of the above.

# Exercise 3

- If \$t0 is holding 17, \$t1 is holding 8, what will be the value stored in \$t2 after the following instructions?

    andi \$t0, \$t0, 3

    beq \$t0, \$0, L1

    addi \$t0, \$t0, 1

L1: add \$t2, \$t0, \$t1

(a) 10.

(b) 8.

(c) 2.

(d) None of the above.

# Exercise 4

- Assume A is an integer array with 10 elements storing 0,1,2,3,4,5,6,7,8,9. Assume the starting address of A is in \$s0 and \$t0 is holding 3. After the running the following code, what will be the content of \$t0?

```
sll $t0, $t0, 3
```

```
add $t0, $s0, $t0
```

```
lw $t0, 0($t0)
```

```
srl $t0, $t0, 1
```

- (a) 3
- (b) 1
- (c) 0
- (d) None of the above.

# In Class Exercise

- If-Else