MIPS coding

Review

- Shifting
 - Shift Left Logical (sll)
 - Shift Right Logical (srl)
 - Moves all of the bits to the left/right and fills in gap with 0's
 - For most cases, equivalent to multiplying/dividing by 2ⁿ where n is the number of bits being shifted
 - Be careful of overflow and using srl on negative numbers

Review

- Branching
 - Branch If Equal (beq)
 - Branch If Not Equal (bne)
 - Jump (j)
 - Changes point of execution:
 - Conditionally only if clause is true for beq/bne (otherwise, the next instruction is the one executed)
 - Unconditionally for j
 - Used to make if statements and loops in higher level languages

In Class Exercise

• Write the MIPS assembly code for the following C code segment:

```
If (A[1] < A[2]) {
A[0] = A[1] & 5;
}
else {
A[0] = A[2] & 5;
}
```

- Assume the starting address of array A is stored in \$s0. Use only the instructions
- covered in class, i.e. add, addi, sub, or, ori, and, andi, xor, xori, nor, lw, sw, srl, sll, beq,
- bne, j.

In Class Exercise – C Code

In Class Exercise – Set up

```
ori $t0, $zero, 5
lw $t1, 4($s0)
lw $t2, 8($s0)
If (A[1] < A[2]) {
         A[0] = A[1] & 5;
}
else {
         A[0] = A[2] \& 5;
}
sw $t0, 0($s0)
```

```
# Set up constant used in if# Get value from A[1] and place in $t1# Get value from A[2] and place in $t2
```

In Class Exercise – If Bodies

ori \$t0, \$zero, 5 lw \$t1, 4(\$s0) lw \$t2, 8(\$s0) If **(\$t1 < \$t2**) { and \$t0, \$t1, \$t0 } else { and \$t0 \$t2, \$t0 } sw \$t0, 0(\$s0)

Set up constant used in if# Get value from A[1] and place in \$t1# Get value from A[2] and place in \$t2

And constant and A[1]

And constant and A[2]

In Class Exercise – Change Compare Operator

ori \$t0, \$zero, 5 lw \$t1, 4(\$s0) lw \$t2, 8(\$s0) sub \$t3, \$t1, \$t2 srl \$t3, \$t3, 31 If **(\$t3 != \$zero**) { and \$t0, \$t0, \$t1 } else { and \$t0, \$t0, \$t2 } sw \$t0, 0(\$s0)

Set up constant used in if # Get value from A[1] and place in \$t1 # Get value from A[2] and place in \$t2 # Negative if <, Zero/Positive if >= # Discard everything but the sign bit

And constant and A[1]

And constant and A[2]

In Class Exercise – Change If Statement to BEQ

ori \$t0, \$zero, 5 lw \$t1, 4(\$s0) lw \$t2, 8(\$s0) sub \$t3, \$t1, \$t2 srl \$t3, \$t3, 31 **beq \$t3, \$zero, ELSE** and \$t0, \$t0, \$t1 **j EXIT**

ELSE:

and \$t0, \$t0, \$t2

EXIT:

sw \$t0, 0(\$s0)

Set up constant used in if # Get value from A[1] and place in \$t1 # Get value from A[2] and place in \$t2 # Negative if <, Zero/Positive if >= # Discard everything but the sign bit

And constant and A[1]

Skip over ELSE branch

And constant and A[2]

In Class Exercise – Convert to Exercise Given (<=)

ori \$t0, \$zero, 5 lw \$t1, 4(\$s0) lw \$t2, 8(\$s0) sub \$t3, \$t1, \$t2 bne \$t3, \$zero, REST and \$t0, \$t0, \$t1 j EXIT

REST:

srl \$t3, \$t3, 31 beq \$t3, \$zero, ELSE and \$t0, \$t0, \$t1 j EXIT

ELSE:

and \$t0, \$t0, \$t2

EXIT:

sw \$t0, 0(\$s0)

Set up constant used in if # Get value from A[1] and place in \$t1 # Get value from A[2] and place in \$t2 # Negative if <, Zero/Positive if >= # Skip if the two numbers are not equal # Same as true branch below # Skip over everything else

Discard everything but the sign bit

And constant and A[1]# Skip over ELSE branch

And constant and A[2]

slt, slti

- slt \$t3, \$t1, \$t2
 - set \$t3 to be 1 if \$t1 < \$t2; else clear \$t3
 to be 0.</pre>
 - "Set Less Than."
- slti \$t3, \$t1, 100
 - set \$t3 to be 1 if \$t1 < 100; else clear \$t3
 to be 0.</pre>

Using slt

- slt \$t3, \$t1, \$t2
 beq \$t3, \$zero, ELSE
- andi \$t0, \$t1, 5
- j EXIT

ELSE:

andi \$t0, \$t2, 5 EXIT:

Complete MIPS code

- The text segment in the source code usually starts with
 - .text .globl main

main:

where ``main'' is the label associated with the address of the first instruction of the code.

• And the code usually ends with

li \$v0,10 # telling the simulator to stop
syscall

• Comment with `#'

In Class Exercise

.text

.globl MAIN

MAIN:

	syscall	# Exits the program
	li \$v0, 10	# Sets the syscall operation
	sw \$t0, 0(\$s0)	# Store \$t0 to A[0]
EXIT:		
	and \$t0, \$t0, \$t2	# And constant and A[2]
ELSE:		
	j EXIT	# Skip over ELSE branch
	and \$t0, \$t0, \$t1	# And constant and A[1]
	beq \$t3, \$zero, ELSE	
	srl \$t3, \$t3, 31	# Discard everything but the sign bit
REST:		
	j EXIT	# Skip over everything else
	and \$t0, \$t0, \$t1	# Same as true branch below
	bne \$t3, \$zero, REST	# Skip if the two numbers are not equal
	sub \$t3, \$t1, \$t2	<pre># Negative if <, Zero/Positive if >=</pre>
	lw \$t2, 8(\$s0)	# Get value from A[2] and place in \$t2
	lw \$t1, 4(\$s0)	# Get value from A[1] and place in \$t1
	ori \$t0, \$zero, 5	# Set up constant used in if

SPIM

- Run codes with **SPIM**. SPIM is a simulator.
 - Use any editor to write the source file, save it as an .asm file.
 - Run SPIM, load the source file.
 - F10 to step through the code. Monitor how the registers change.
 - F5 to run the code
 - Can set breakpoints for debugging
- SPIM can be downloaded at

http://sourceforge.net/projects/spimsimulator/files/

• Lots of good references online, like

https://www.cs.tcd.ie/~waldroj/itral/spim_ref.html

Working with the simulator

- Can check
 - How the program runs
 - How the instructions are encoded, addressed
 - How to monitor the change of the registers
 - Later, how the memory is used to store data

Some Comments

- Being able to write if-else, we can have all other fancy things like for loop, while loop....
- That is why we do not have an instruction for the for loop or while loop, but we build it from the if-else.

Compiling a while loop in C

How to translate the following to MIPS assembly?
 while (save[i] == k)

 i += 1;

- We first translate into a C program using if and goto

```
Loop: if (save[i] != k) goto Exit;
i = i + 1;
goto Loop;
Exit:
```

Compiling a while loop in C

Assume that i and k correspond to registers
 \$s3 and \$s5 and starting address of array save
 is in \$s6

while (save[i] == k) i += 1;

Compiling a while loop in C

Assume that i and k correspond to registers
 \$s3 and \$s5 and starting address of array save
 is in \$s6

Exit:

While Loop

 How many instructions will be executed for the following array save?

10, 10, 10, 10, 10, 10, 10, 10, 10, 0

- Assume that k = 10 and i = 0 initially

- (6 loop lines * 9 loops) + 4 lines in last iteration $\frac{-}{5/28/2013}$ = 58 lines week04-3.ppt

Optimized

sll \$t1, \$s3, 2 # Temp reg \$t1 = 4 * i add \$t1, \$t1, \$s6 # \$t1 = address of save[i] lw \$t0, 0(\$t1) # Temp reg \$t0 = save[i] bne \$t0, \$s5, Exit # go to Exit if save[i] ≠ k Loop: addi \$s3, \$s3, 1 # i = i + 1 addi \$t1, \$t1, 4 # \$t1 = address of save[i] lw \$t0, 0(\$t1) # Temp reg \$t0 = save[i] beq \$t0, \$s5, Loop # go to Loop if save[i] = k Exit:

- How many instructions now?
 - Assume k = 10 and i = 0 initially

10, 10, 10, 10, 10, 10, 10, 10, 10, 0

- 4 preloop lines + (4 loop lines * 9 loop iterations) + 4 lines in last iteration
- = 44 lines

The loop code

```
.data
save:.word 10, 10, 10, 10, 10, 11, 12,
    .text
    .globl main
main:
    li $s3, 0
    li $s5, 10
    la $s6, save
Loop:
    sll $t1, $s3, 2
    add $t1, $t1, $s6
    lw $t0, 0($t1)
    bne $t0, $s5, Exit
    addi $s3, $s3, 1
    j Loop
Exit:
done:
    li $v0, 10 # these two lines are to tell the simulator to stop
```

syscall

Data segment and code segment

- The code has a **data** segment and a **code (text)** segment.
- The beginning of the data segment in the assembly source code is indicated as

.data

and followed by several declarations such as

meaning an array of words whose starting address is associated with label ``A.''

- Several notes:
 - It will allocate continuous spaces in the memory for the data
 - .word means everything is 4 bytes
 - save: is a label associated with the address of the first byte allocated. Like the label for the instructions, label for an address is also an address.